

On $(K_q; k)$ -Stable Graphs

by Matthew Ridge

Under the direction of
Prof. John Caughman

With second reader
Prof. Paul Latiolais

On $(K_q; k)$ -Stable Graphs

An article by Andrzej Żak, first published online
the 15th of October 2012.

Introduction

- Definitions
- General Bounds
- Complete Graphs, K_q
- Future Work
- Conclusion

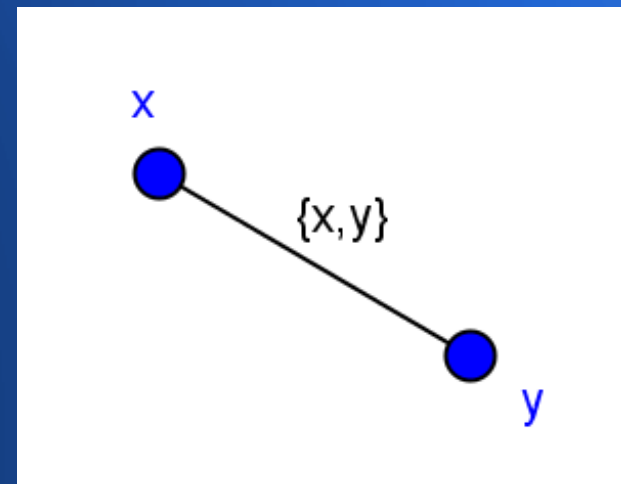
Definitions

- What is a graph?
 - Order of a graph.
 - Size of a graph.
- Graph Isomorphism.
- Cycles, C_n
- Complete Graphs, K_n
- Vertex Stable Graph.
 - Minimum vertex stable graph.
 - Size of a minimum vertex stable graph.

What is a graph?

A (simple) graph G consists of a *vertex* set $V(G)$ and an *edge* set $E(G)$ where $E(G)$ is a set of 2-element subsets of $V(G)$. When $\{x, y\} \in E(G)$ we write $x \sim y$. A vertex that does not belong to any edge is called an *isolated vertex*.

Let x, y be vertices of G , then if x and y share an edge we denote that edge as $\{x, y\}$.



What is a graph?

- Order of a Graph: Let G be a graph, then the *order* of G is defined as the number of vertices it has, denoted:

$$|G| := |V(G)|.$$

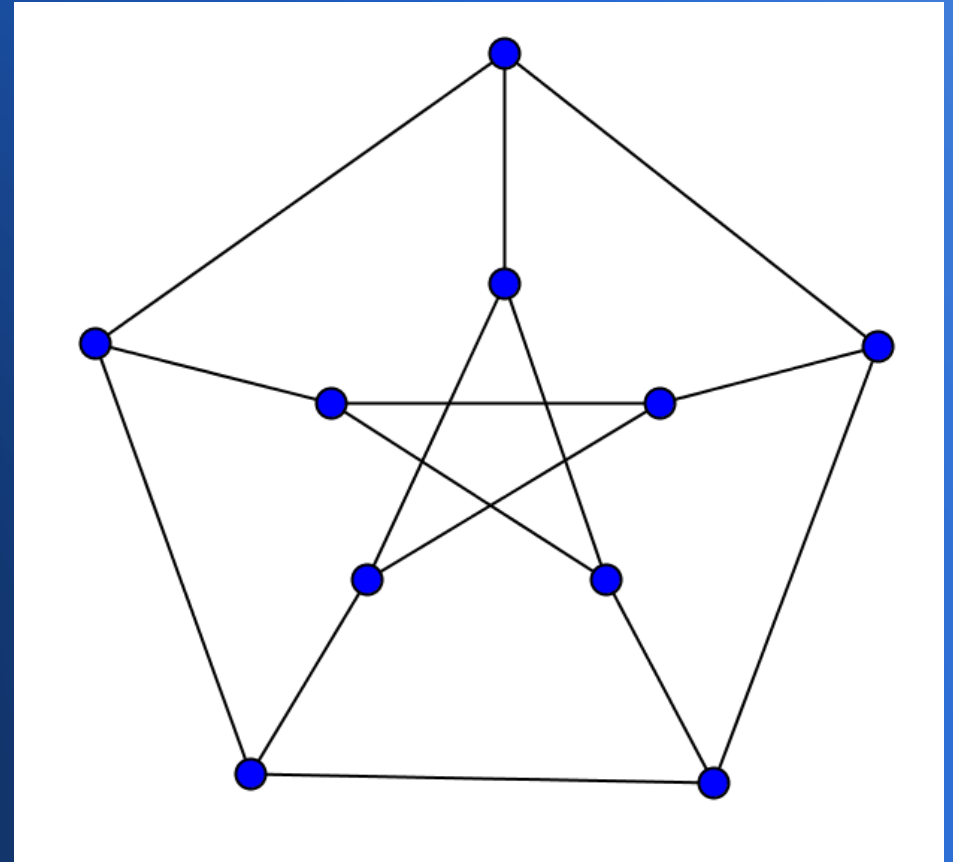
- Size of a Graph: Let G be a graph, then the *size* of G is defined as the number of edges it has, denoted:

$$\|G\| := |E(G)|.$$

Example

Call the following graph G .

- What is the Order of this graph?
- What is the size of the graph?



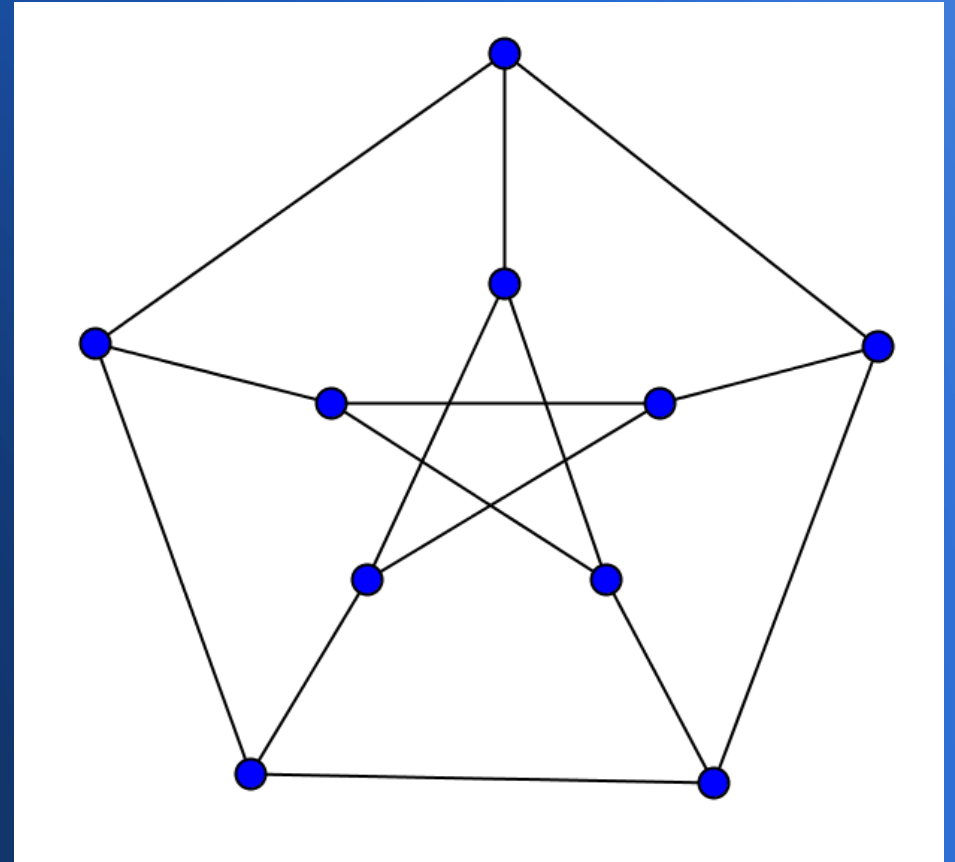
Example

- What is the Order of this graph?

$$|G|=10$$

- What is the size of the graph?

$$\|G\|=15$$



Graph Isomorphism

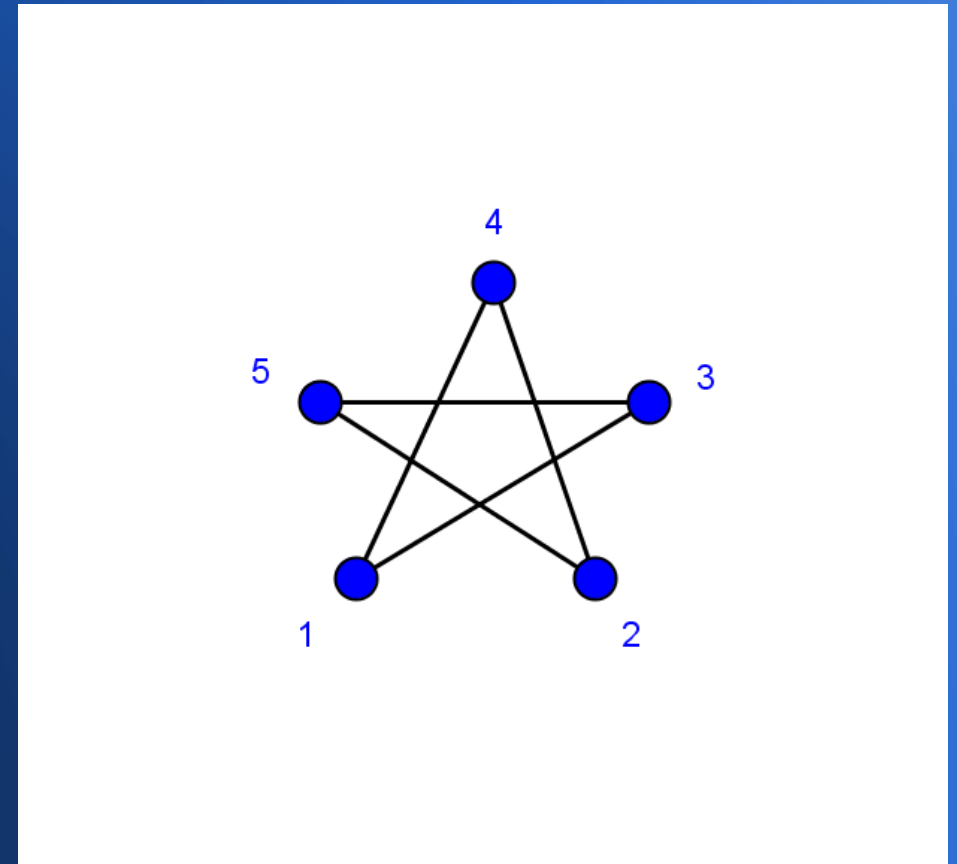
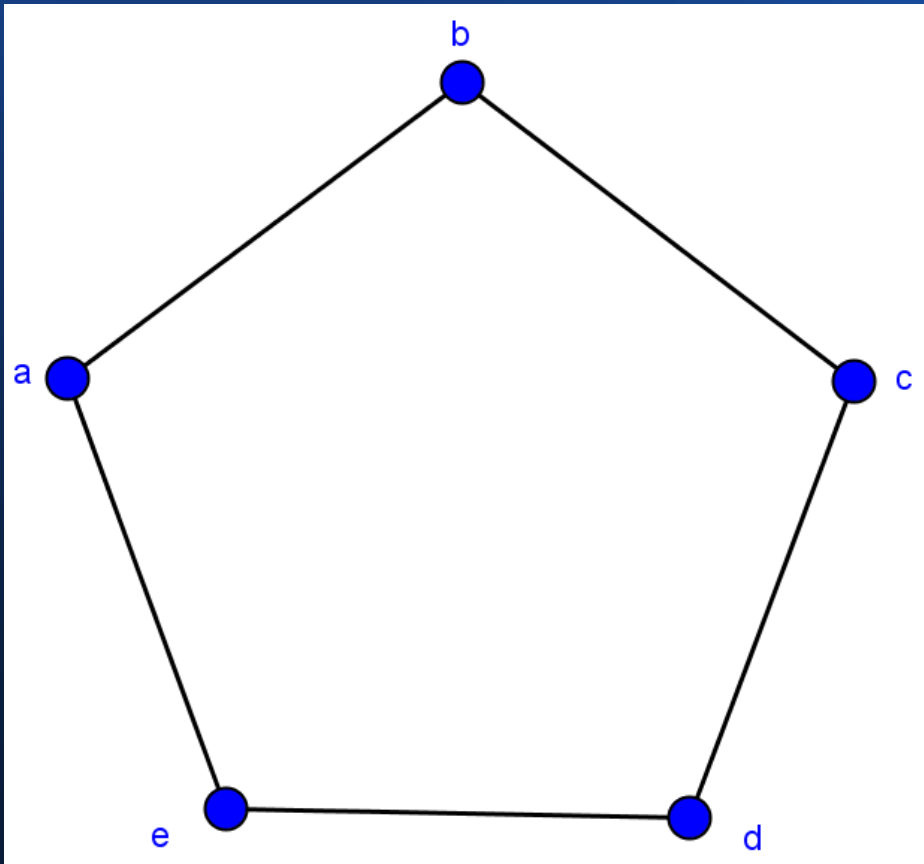
Two graphs G and H are *isomorphic* if there is a bijection

$$f : V(G) \rightarrow V(H)$$

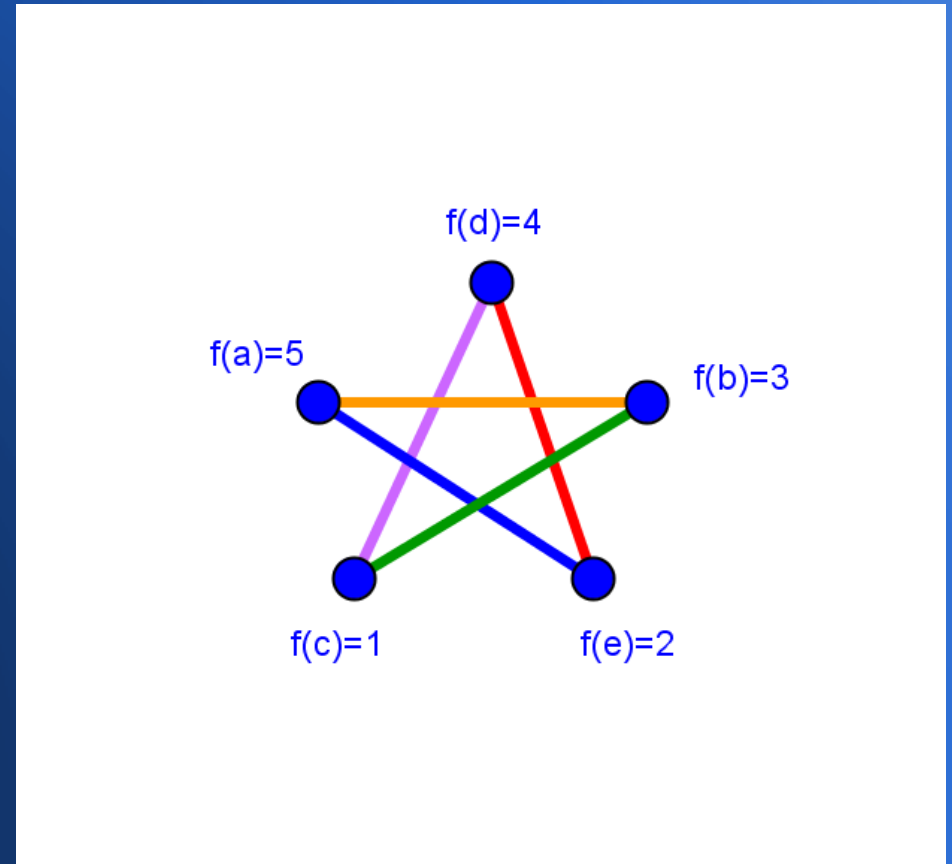
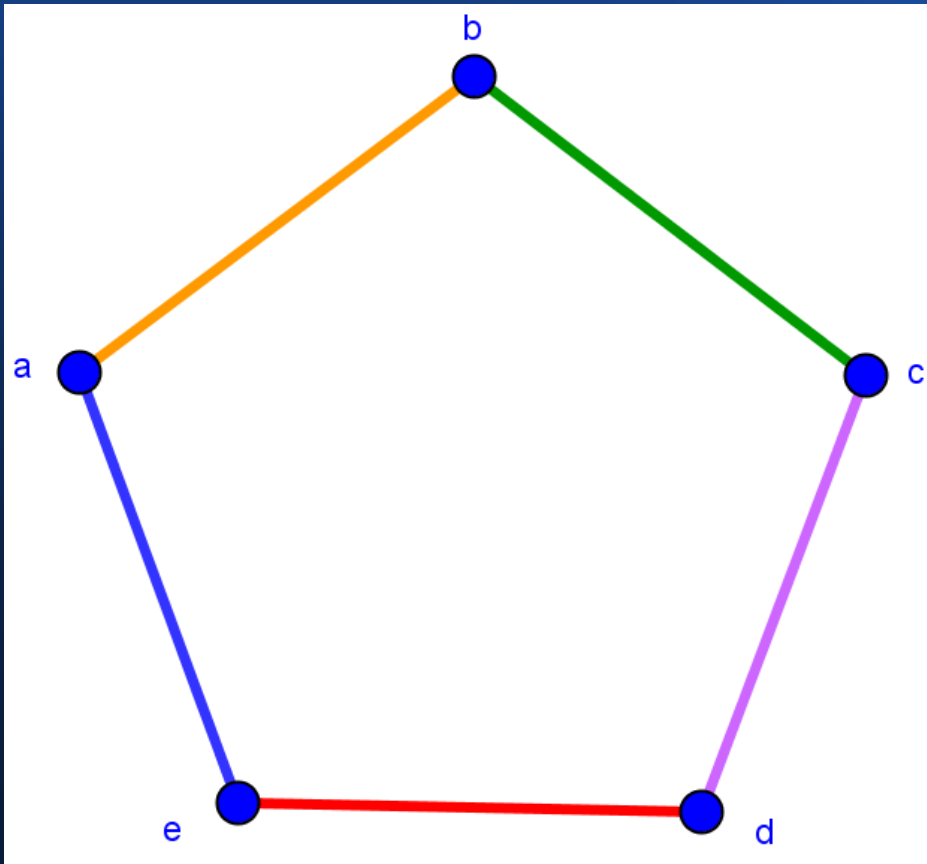
such that

$$u \sim v \text{ iff } f(u) \sim f(v).$$

Example



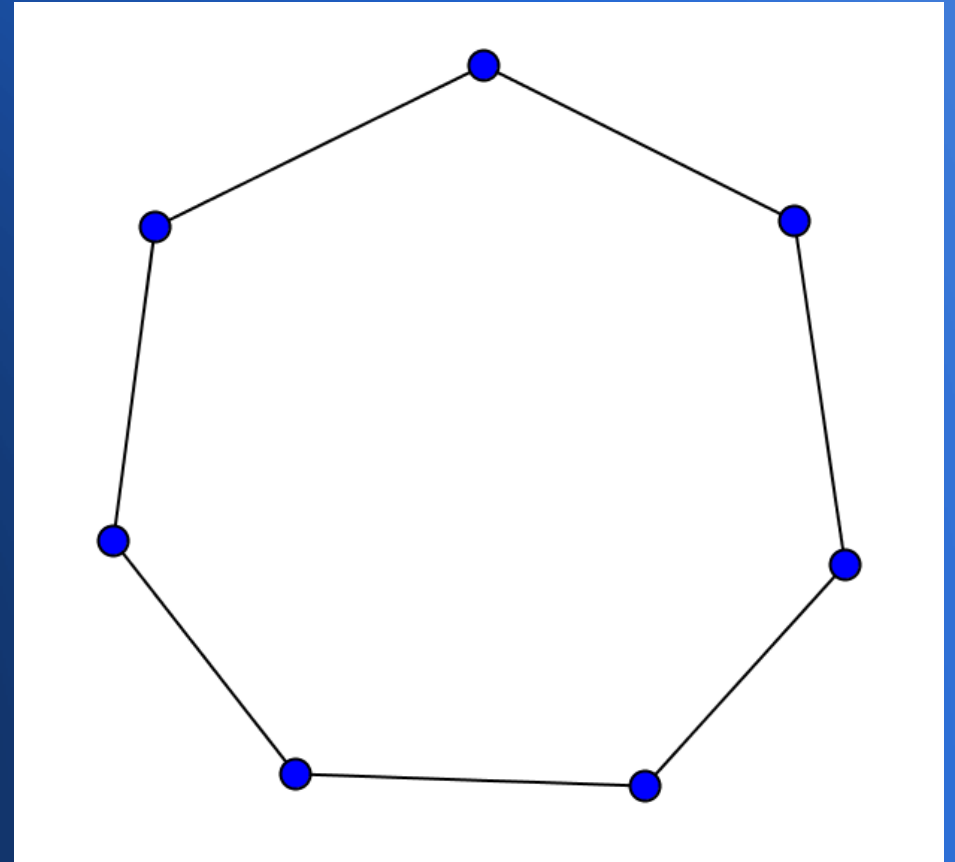
Example



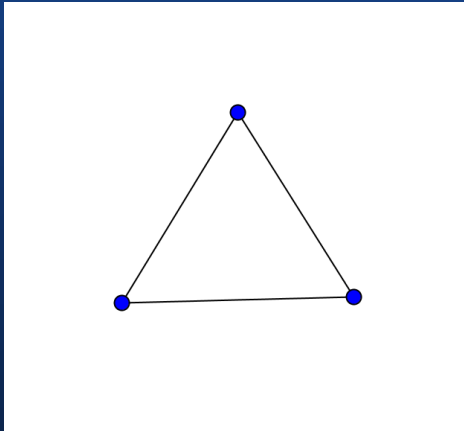
Cycle graphs, C_n

A cycle, more commonly called a closed walk, consists of a sequence of vertices starting and ending at the same vertex, with each two consecutive vertices in the sequence adjacent to each other in the graph.

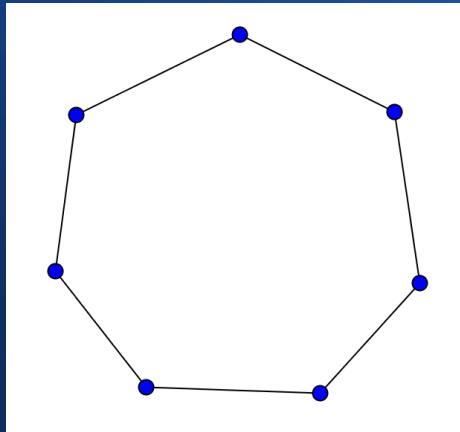
Example: C_7



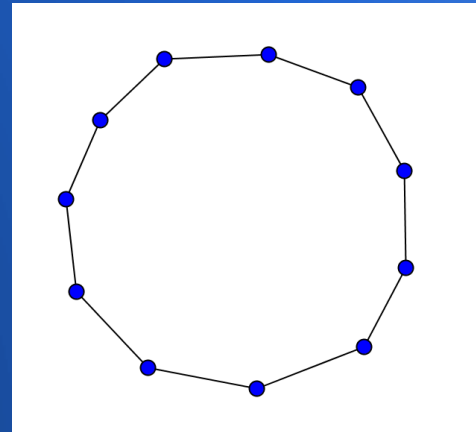
Cycle graphs, C_n



C_3



C_7

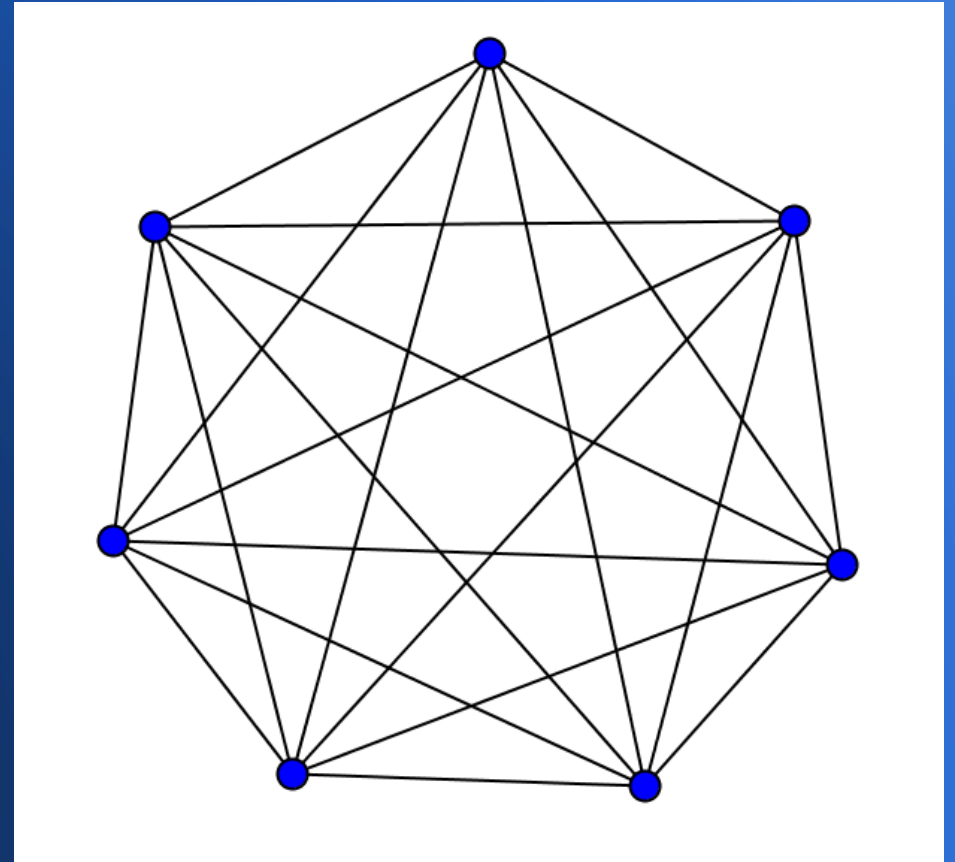


C_{11}

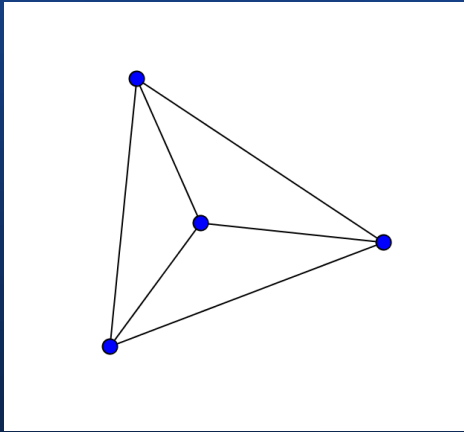
Complete Graphs, K_q

A complete graph is a simple undirected graph in which every pair of distinct vertices is connected by a unique edge.

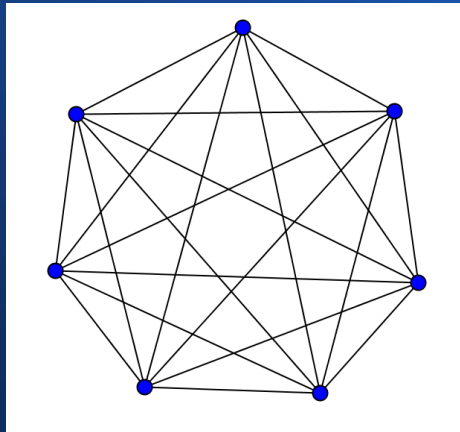
Example: K_7



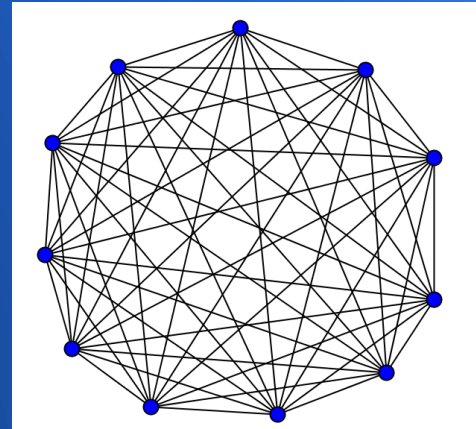
Complete Graphs, K_q



K_4



K_7



K_{11}

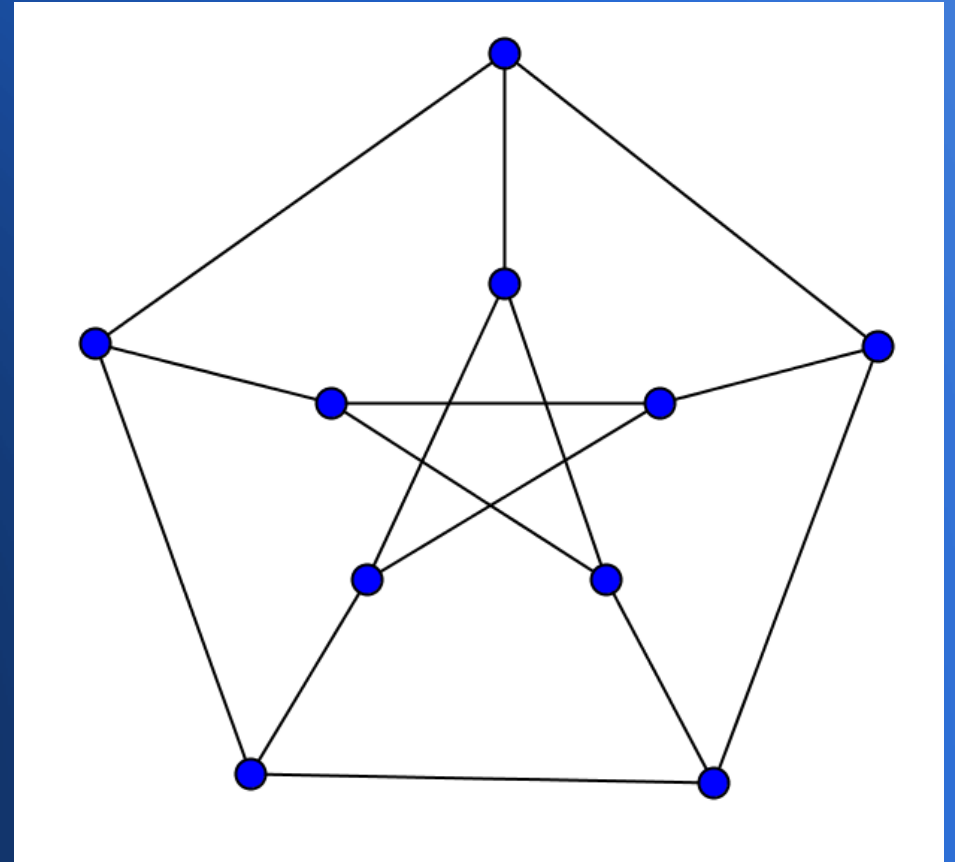
Vertex Stable

A graph G is called *$(H;k)$ -vertex stable* or *$(H;k)$ -stable* if G contains a subgraph isomorphic to H even after removing any k of its vertices.

Example: Consider the Petersen Graph and we're going to check its stability with $H=C_5$.

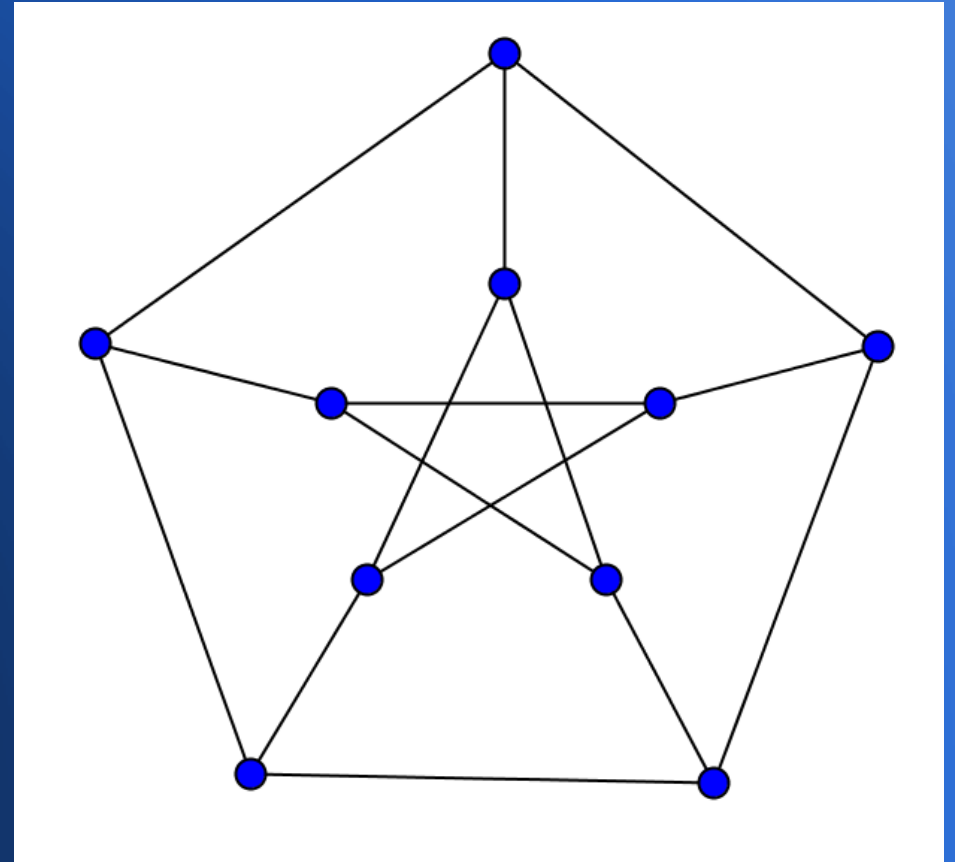
Example

Is the Petersen Graph
($C_5; 0$)-stable?



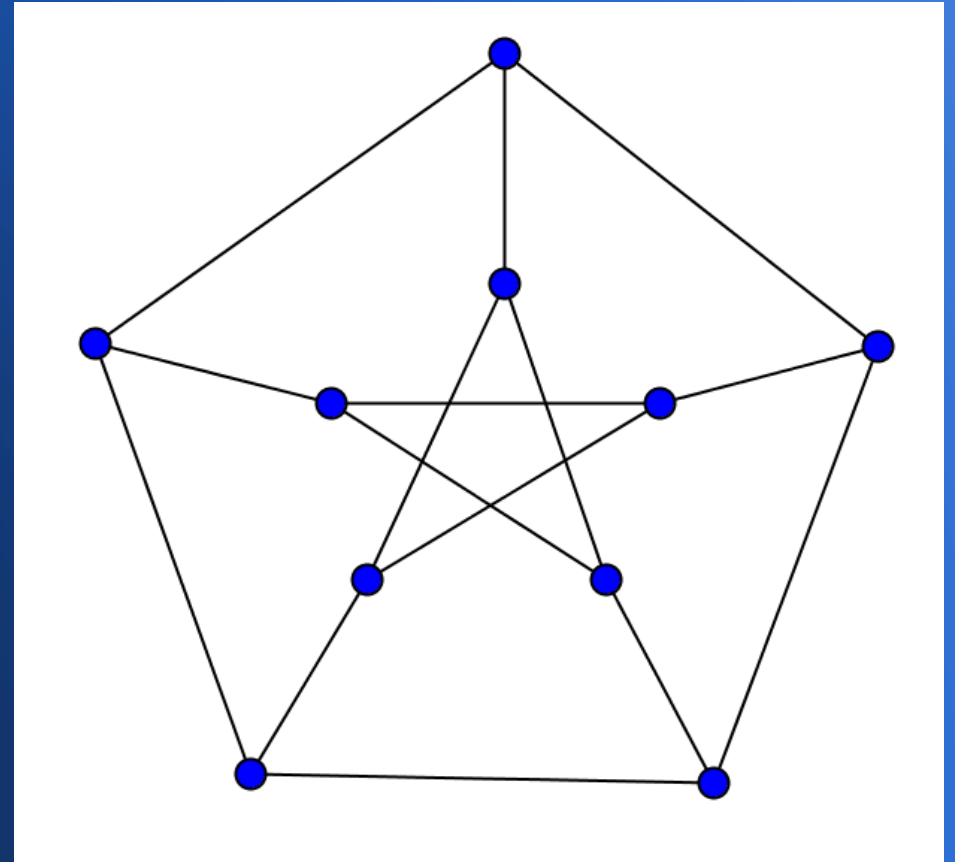
Example

Is the Petersen Graph
 $(C_5, 1)$ -stable?

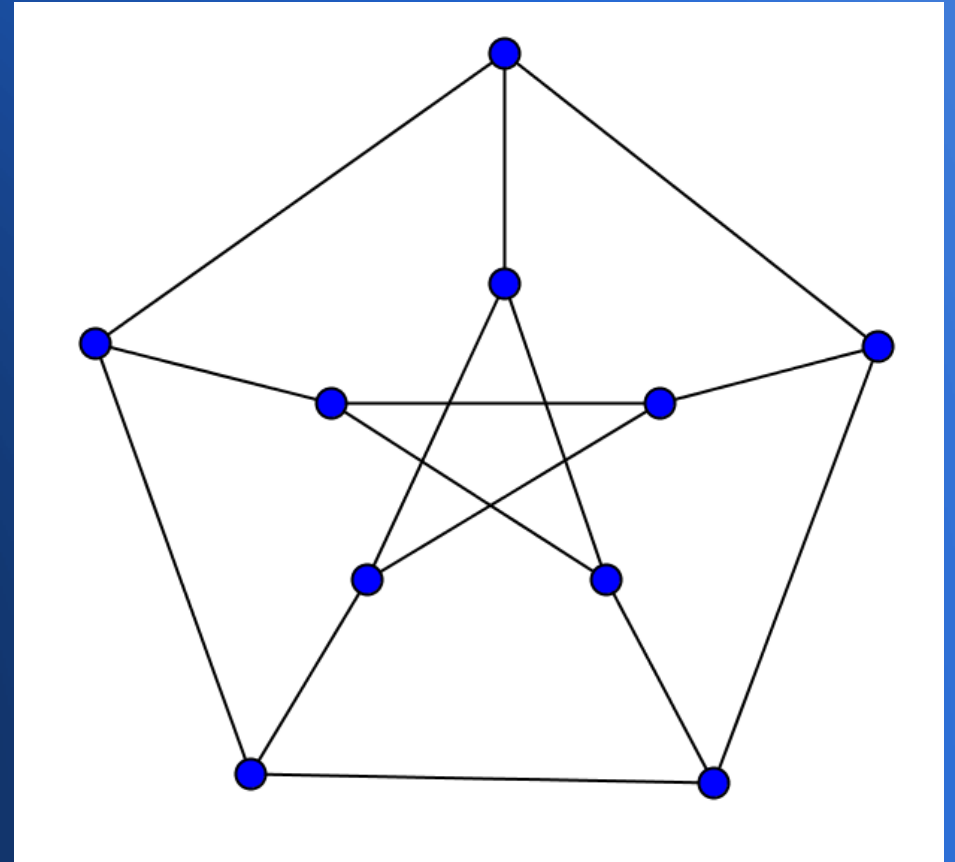
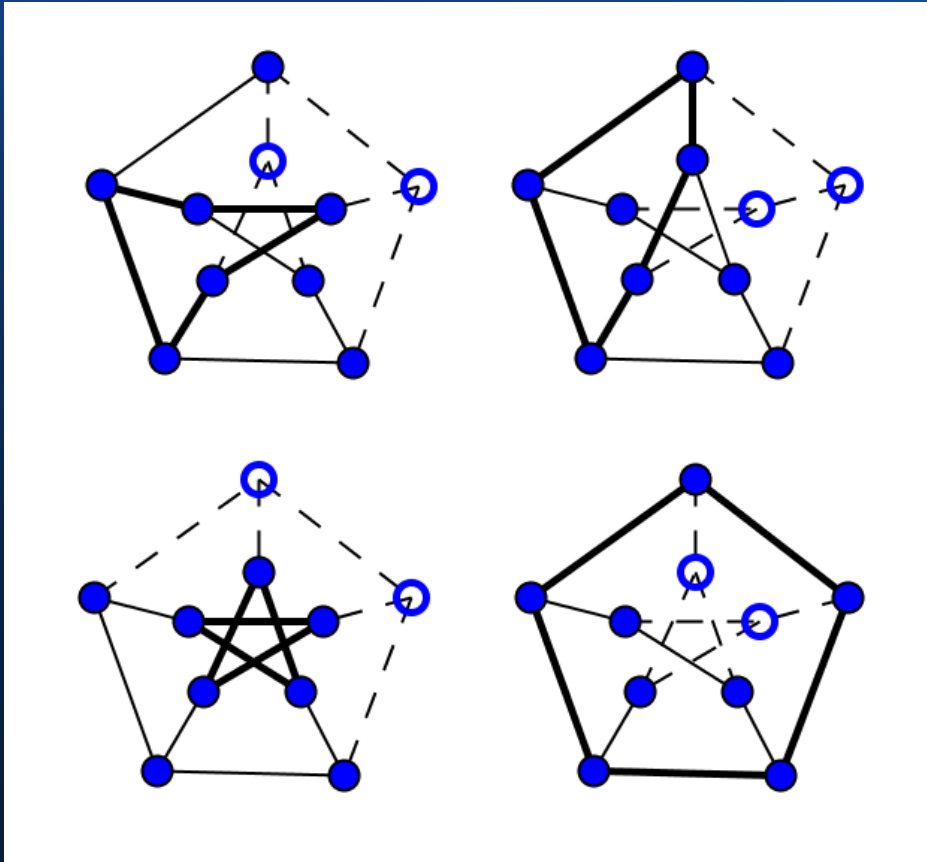


Example

Is the Petersen Graph
 $(C_5; 2)$ -stable?

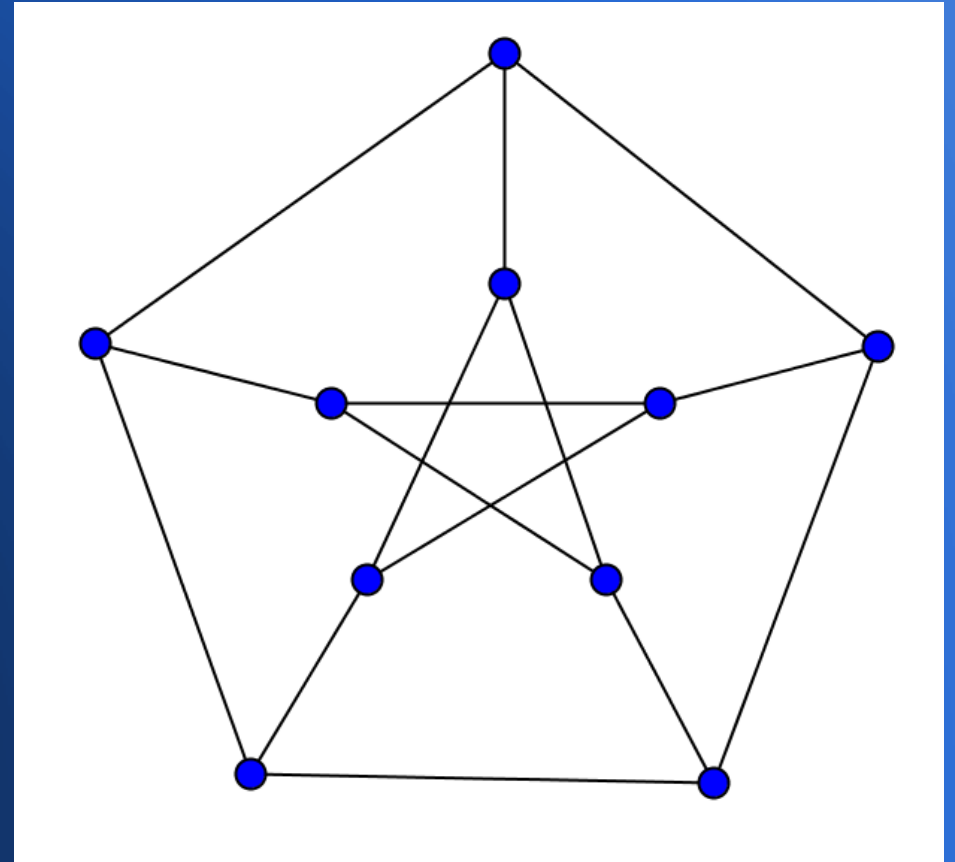


Example

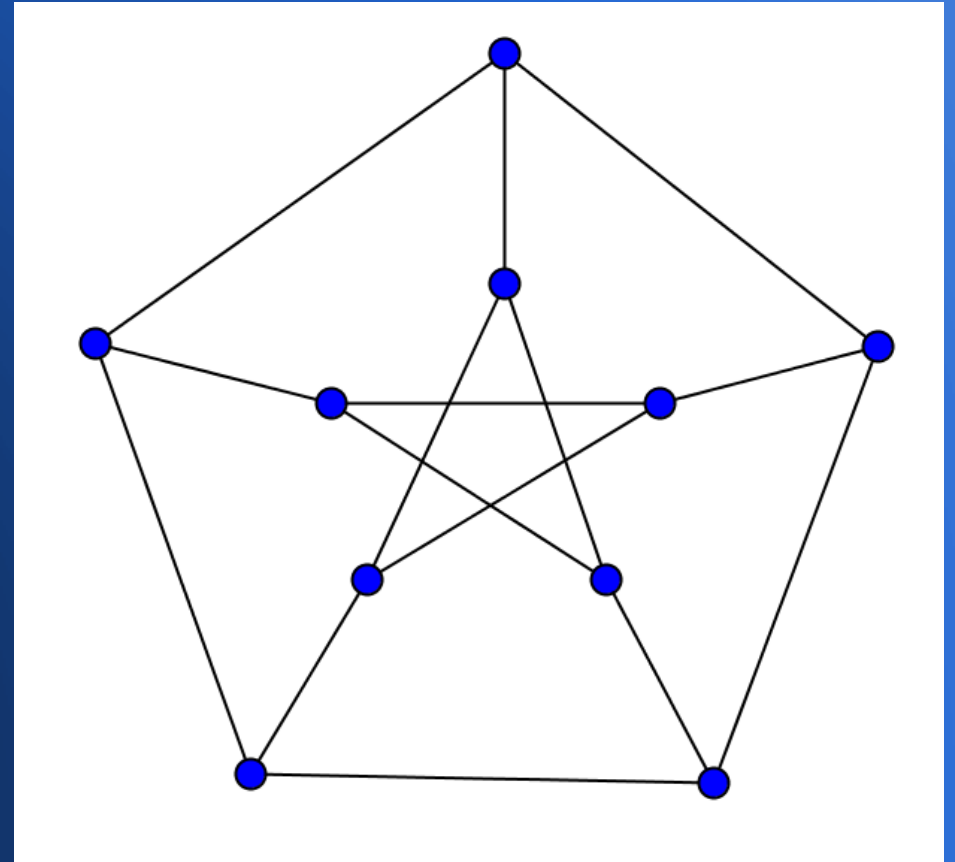
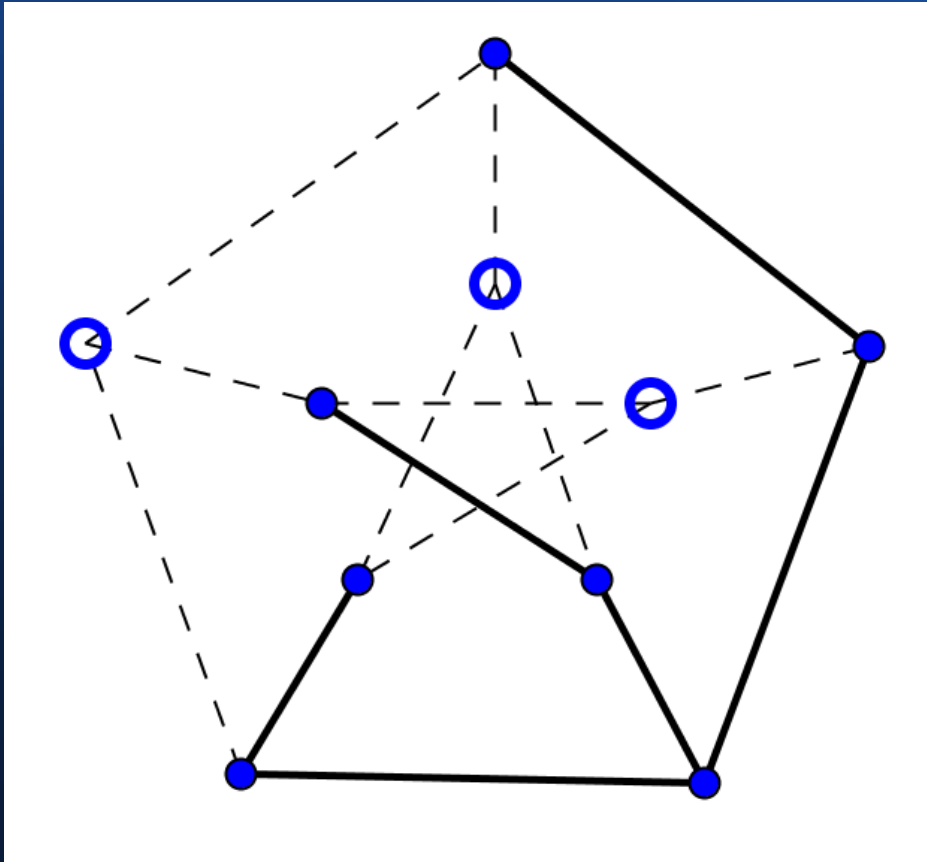


Example

Is the Petersen Graph
 $(C_5, 3)$ -stable?



Example



minimum $(H;k)$ -stable graph

If G has the minimum size of any $(H;k)$ -stable graph, then

$$\|G\| := \text{stab}(H ; k),$$

and we refer to G as a minimum $(H;k)$ -stable graph.

Further assumptions

We will not consider isolated vertices in any of the graphs in question.

Why? After adding to or removing from an $(H;k)$ -stable graph any number of isolated vertices, we still have an $(H;k)$ -stable graph with the same size.

General Bounds

The goal is to find a lower bound on the size of our stable graph, G , for selected graph H and arbitrarily chosen k .

Major problem to address:

We need to be able to remove *any* set of k vertices without losing a subgraph of H .

Theorem

Let δ_H be the minimum degree of H . If G is a minimum $(H;k)$ -stable graph, then

$$|G| - \delta_H \sum_{v \in V(G)} \frac{1}{d_G(v) + 1} \geq k + 1.$$

Moreover, if G is not a union of cliques, then the inequality is strict.

(This is the light at the end of the tunnel.)

Lemma

If G is a minimum $(H;k)$ -stable graph, then every vertex and every edge of G belongs to some subgraph of G isomorphic to H .

Proof: Roughly, imagine that you have a minimum $(H;k)$ -stable graph where there is a vertex or edge that doesn't belong to a subgraph of G isomorphic to H . If that edge or vertex is never used in any of the isomorphisms, then its removal doesn't affect the stability of G . Hence G was not minimum! ■

Proposition

Let δ_H be the minimum degree of a graph H . Then in any minimum $(H;k)$ -stable graph G ,

$$d_G(v) \geq \delta_H$$

for each vertex $v \in V(G)$.

Proof: Follows from the previous lemma.

The Ordering...

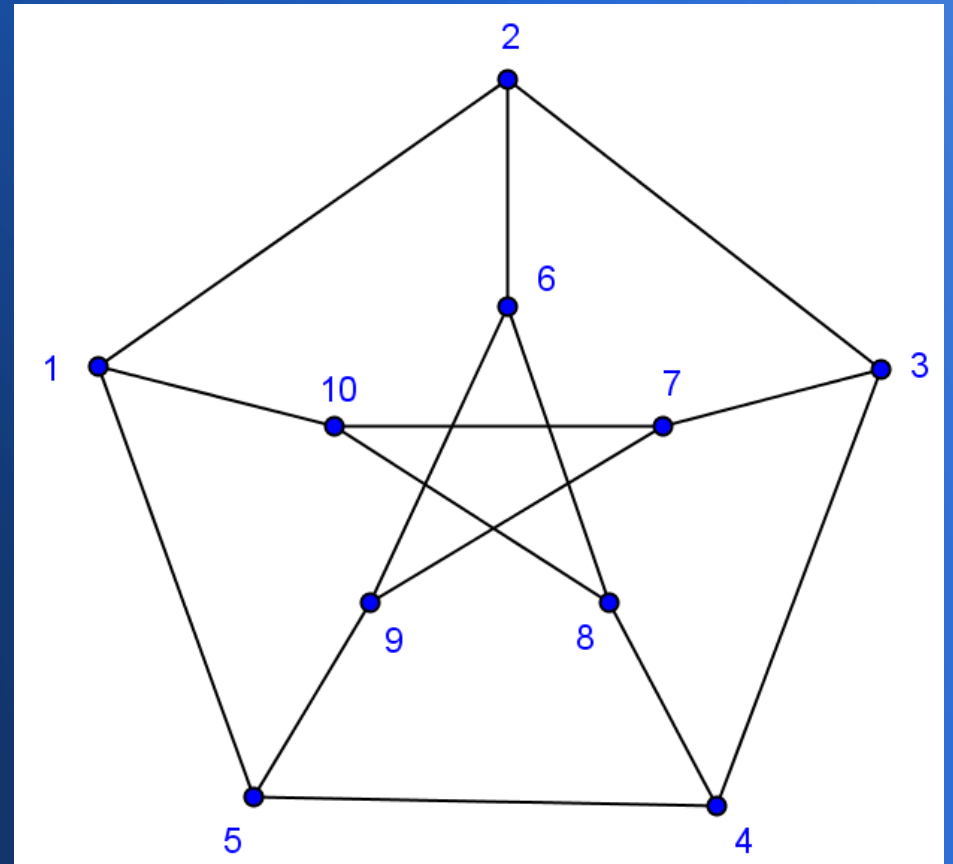
Fix any ordering σ of $V(G)$. Let $deg_{\sigma}(v)$ denote the number of neighbors of v that are on the left of v in ordering σ . Then let S_{σ} denote the set of all vertices with $deg_{\sigma}(v) \leq \delta_H - 1$.

If remove from G the set of vertices in $V(G) \setminus S_{\sigma}$, the claim is that this will induce a subgraph on G that will contain no copies of H .

What?

Example

Consider a labeling of our previous graph, the Petersen Graph. Further, suppose that H is again C_5 .



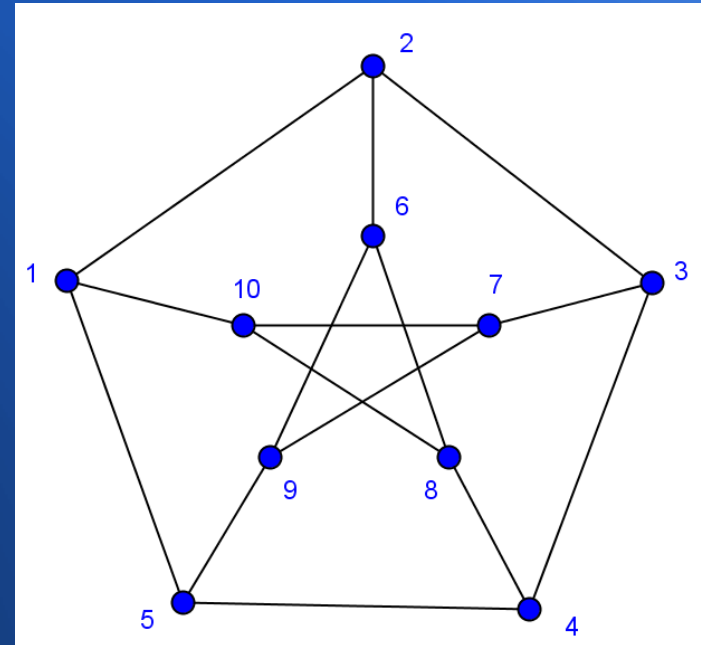
Example

Observe that

$$\begin{aligned} S_\sigma &:= \{v_i \mid \deg_\sigma(v_i) \leq 1\} \\ &= \{v_1, v_2, v_3, v_4, v_6, v_7\} \end{aligned}$$

Hence

$$V(G) \setminus S_\sigma := \{v_5, v_8, v_9, v_{10}\}$$



n	1	2	3	4	5	6	7	8	9	10
$\deg_\sigma(v_n)$	0	1	1	1	2	1	1	2	3	3

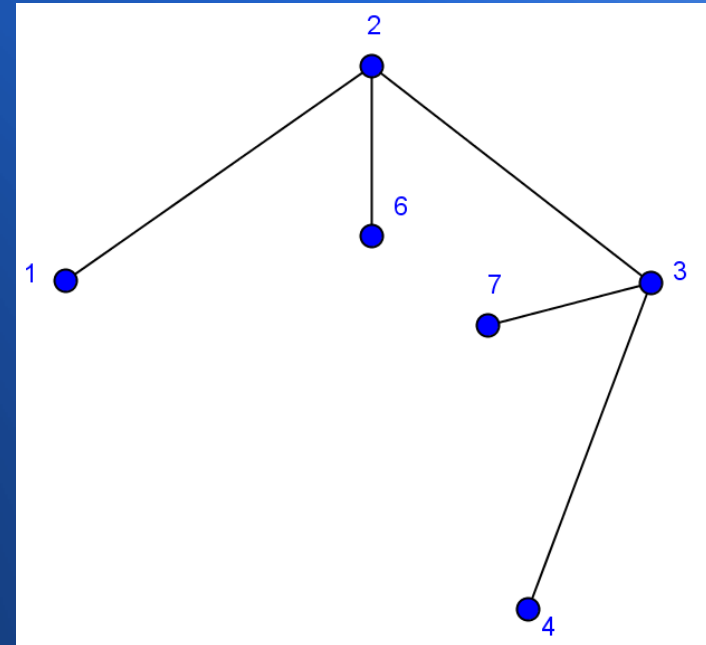
Example

Observe that

$$\begin{aligned} S_\sigma &:= \{v_i \mid \deg_\sigma(v_i) \leq 1\} \\ &= \{v_1, v_2, v_3, v_4, v_6, v_7\} \end{aligned}$$

Hence

$$V(G) \setminus S_\sigma := \{v_5, v_8, v_9, v_{10}\}$$



n	1	2	3	4	5	6	7	8	9	10
$\deg_\sigma(v_n)$	0	1	1	1	2	1	1	2	3	3

So how does it work?

In general, with any ordering we destroy all copies of H by consecutively eliminating all vertices of S_σ . Each vertex in S_σ has left degree $\leq \delta_H - 1$ and thus cannot be the rightmost vertex of any copy of H in the induced subgraph.

From this we get the following:

$$|G| - |S_\sigma| \geq k + 1$$

What next?

We need to find a way to approximate the size of S_σ . To do this we will need...

Probability and Expected Values.

Lets count some stuff!

The story.

Our story begins with letting σ be an ordering on a vertex set of size n ... sitting somewhere in this ordering is an arbitrary vertex v , with degree equal to $d_G(v)$. Also sitting within this ordering is all of v 's neighbors.

— ... — ... — ... — ... — ... — ... —
— ... — ... — ... v ... — ... — ... —
v₁ ... v₂ ... v₃ ... v ... v₄ ... v₅ ... v₆

Choosing spots

The vertex v has $d_G(v)$ neighbors for which each gets a spot in the ordering. Don't forget v needs a spot too! So we get the following:

$$\binom{n}{d_G(v)+1}$$

Choosing spots

Recall that $deg_{\sigma}(v) \leq \delta_H - 1$. We have that if we assign v to any of these first δ_H spots, we satisfy $deg_{\sigma}(v) \leq \delta_H - 1$. Hence we get:

$$\binom{n}{d_G(v)+1} (\delta_H)$$

The rest of the counting is assigning the remaining vertices to their spots and dividing by the total number of permutations on n vertices.

The Probability

$$\begin{aligned} \Pr(\deg_{\sigma}(v) < \delta_H) &= \frac{\binom{n}{d_G(v)+1} (\delta_H)(d_G(v))! (n-d_G(v)-1)!}{n!} \\ &= \frac{\delta_H}{d_G(v)+1} \end{aligned}$$

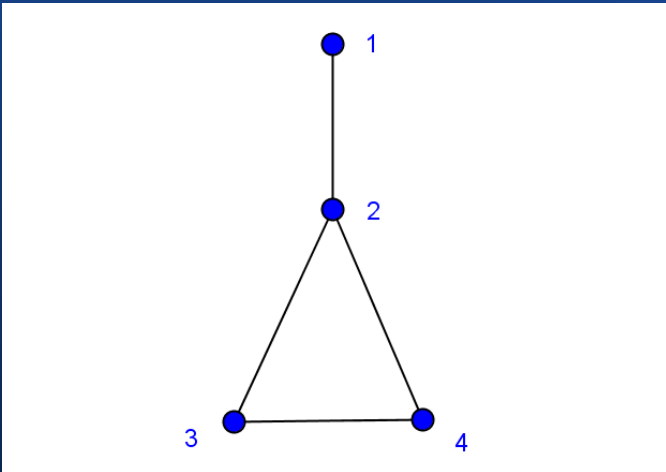
The Expectation

By using a method of expected values we get...

$$Pr(v \in S_\sigma) = \frac{\delta_H}{d_G(v) + 1}$$

$$E(|S_\sigma|) = \sum_{v \in V(G)} \frac{\delta_H}{d_G(v) + 1}$$

Example



1234	2134	3124	4123
1243	2143	3142	4132
1342	2314	3214	4213
1324	2341	3241	4231
1432	2413	3412	4312
1423	2431	3421	4321

Set $\delta_H = 2$.

$$E(|S_\sigma|) = \sum_{v \in V(G)} \frac{\delta_H}{d_G(v) + 1} = \frac{2}{1+1} + \frac{2}{3+1} + \frac{2}{2+1} + \frac{2}{2+1}$$
$$= \frac{24}{24} + \frac{12}{24} + \frac{16}{24} + \frac{16}{24} = \frac{68}{24} \approx 2.8\overline{33}$$

Thus...

$$E(|S_\sigma|) = \sum_{v \in V(G)} \frac{\delta_H}{d_G(v) + 1}$$

Gives us that

$$|G| - |S_\sigma| \geq k + 1$$

$$|G| - \sum_{v \in V(G)} \frac{\delta_H}{d_G(v) + 1} \geq k + 1$$



Corollary

Let H be any graph and let δ_H denote the minimum degree of H . Then

$$\text{stab}(H; k) \geq (k+1) \left(\delta_H + \sqrt{\delta_H(\delta_H - 1)} - \frac{1}{2} \right).$$

What?

Corollary

Roughly, with the use of...

Lemma: The expression

$$\sum_{k=1}^l \frac{1}{x_k} \quad \text{with} \quad \sum_{k=1}^l x_k = r, r \in \mathfrak{R}$$

is minimal if all the x_j are equal. (Uses constrained optimization.)

Example: Set $r=30$,

$$30 = 9 + 8 + 7 + 2 + 2 + 2 \quad \frac{1}{9} + \frac{1}{8} + \frac{1}{7} + \frac{1}{2} + \frac{1}{2} + \frac{1}{2} = \frac{947}{504} \approx 1.879$$

$$30 = 5 + 5 + 5 + 5 + 4 + 6 \quad \frac{1}{5} + \frac{1}{5} + \frac{1}{5} + \frac{1}{5} + \frac{1}{4} + \frac{1}{6} = \frac{73}{60} \approx 1.2167$$

$$30 = 5 + 5 + 5 + 5 + 5 + 5 \quad \frac{1}{5} + \frac{1}{5} + \frac{1}{5} + \frac{1}{5} + \frac{1}{5} + \frac{1}{5} = \frac{6}{5} = 1.2$$

Corollary

Using this fact, note that if the average degree of G is defined:

$d_G = \frac{2\|G\|}{|G|}$ we get: $\|G\| = \frac{d_G |G|}{2}$. Then

$$\sum_{v \in V(G)} \frac{1}{d_G(v)+1} \geq \sum_{v \in V(G)} \frac{1}{d_G+1} = \frac{|G|}{d_G+1}.$$

Rearranging the results from our first theorem we get

$$|G| \geq k+1 + \sum_{v \in V(G)} \frac{\delta_H}{d_G(v)+1} \geq k+1 + \frac{|G|}{d_G+1}.$$

Putting all of this together we and applying a bunch of algebra...

$$|G| \geq (k+1) \frac{d_G+1}{d_G+1-\delta_H}.$$

Corollary

...and using the quadratic formula on a characteristic polynomial and then applying the first derivative test to verify the minimum point. It follows that...

$$\|G\| = \left(\frac{d_G}{2} \right) |G| \geq \frac{k+1}{2} \cdot \frac{d_G(d_G+1)}{d_G+1-\delta_H} \geq (k+1) \left(\delta_H + \sqrt{\delta_H(\delta_H-1)} - \frac{1}{2} \right).$$

Since G was assumed to be minimal we get that...

$$\text{stab}(H; k) \geq (k+1) \left(\delta_H + \sqrt{\delta_H(\delta_H-1)} - \frac{1}{2} \right).$$

Complete Graphs, K_q

Nice thing about complete graphs: All the vertices have the same degree... so if we set H to be K_q , then $\delta_H = q - 1$.

Couple this fact with our previous theorem, corollary, and TONS of algebra, we get...

Theorem

Let G be a $(K_q; k)$ -stable graph, $q \geq 2$ and $k \geq 0$ then

$$\|G\| \geq (2q-3)(k+1)$$

with equality if and only if G is a disjoint union of cliques K_{2q-3} and K_{2q-2} .

The equality is simply a nice consequence from selecting cliques of consecutive sizes. Simple algebra proves this... A lot of simple algebra.

The coin problem...

is a mathematical problem that asks for the largest monetary amount that cannot be obtained using only 2 coins of specified denominations (relatively prime).

In other words, there should exist a positive integer such that for any integer greater than it, there exists a linear combination of the two denominations to express the value.

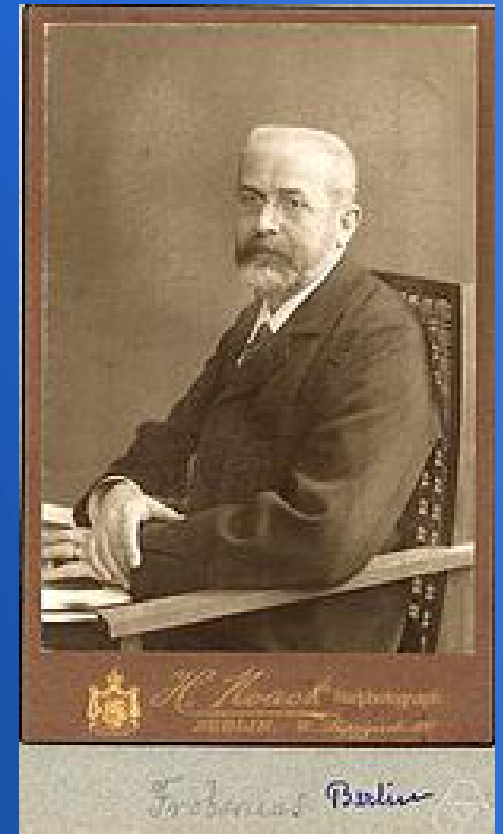
Frobenius Numbers

The largest nonnegative integer that can't be expressed as a linear combination of the two relatively prime denominations is called the Frobenius Number.

Idea by Georg Frobenius.

The following by James Joseph Sylvester:

Given positive integers m, n where $\gcd(m, n) = 1$, the Frobenius number $m(a) + n(b) = K$ is given by $mn - m - n$.



Georg Frobenius

Example

$$4(1) + 5(0) = 4$$

$$4(0) + 5(1) = 5$$

$$4(2) + 5(0) = 8$$

$$4(1) + 5(1) = 9$$

$$4(0) + 5(2) = 10$$

$$4(3) + 5(0) = 12$$

$$4(2) + 5(1) = 13$$

$$4(1) + 5(2) = 14$$

$$4(0) + 5(3) = 15$$

⋮

Confirming with the closed formula:

$$m = 4, n = 5$$

$$mn - m - n$$

$$4 \cdot 5 - 4 - 5 = 11$$

Frobenius number for 4,5 is 11!

McNugget Problem is a “famous” example of the use for Frobenius Numbers.

Theorem

Let $q \geq 2$, $k \geq 0$ be non-negative integers. Then

$$\text{stab}(K_q; k) \geq (2q-3)(k-1),$$

with equality if and only if $k = a(q-2) + b(q-1) - 1$ for some non-negative integers a, b . In particular,

$$\text{stab}(K_q; k) = (2q-3)(k-1), \text{ for } k \geq (q-3)(q-2) - 1.$$

Furthermore, if G is $(K_q; k)$ -stable graph with

$$\|G\| = (2q-3)(k+1)$$

then G is a disjoint union of cliques K_{2q-3} and K_{2q-2} .

Future Work

- Things to think about:
 - Uniqueness of minimal stable graph?
 - Families of graphs that guarantee certain stability. For example...
- Can prove:
 - C_{kn+k+1} is $(P_n; k)$ -stable.
- Conjecture (Fermat-style):
 - C_{kn+k+1} is the unique minimal $(P_n; k)$ -stable graph.
 - $K_{q,r,s}$ is $(C_{2n+1}; k)$ -stable.
 - $K_{q,r}$ is $(C_{2n}; k)$ -stable.

References

- N. Alon and J. Spencer, The Probabilistic Method, John Wiley, New York, NY, 2nd edition, 2000.
- A. Dudek, A. Szymański, and M. Zwonek, $(H; k)$ stable graphs with minimum size, Discuss Math Graph Theory 28 (1) (2008), 137–149.
- S. Cichacz, A. Görlich, M. Zwonek, and A. Żak, On $(C_n; k)$ stable graphs, Electron J Combin 18 (1) (2011), #P205.
- A. Żak, On $(K_q; k)$ -Stable Graphs, J. Graph Theory 74(2) (2013), 216-221.

The End

Thank you!

Special Thanks:

My committee:
John Caughman
and Paul Latiolais

My wife, Temris
Mark and Janet

